

# Efficient Modeling of Half-reified global constraints

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European Research Council  
Established by the European Commission

**Tuples**

TRUSTWORTHY AI

# Modeling systems



Modeling system



# Modeling with global constraints



*AllDifferent(x, y, z)*

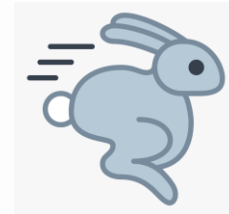
Modeling system



*AllDifferent(x, y, z)*



+ Easy for modeler



+ Fast propagation!

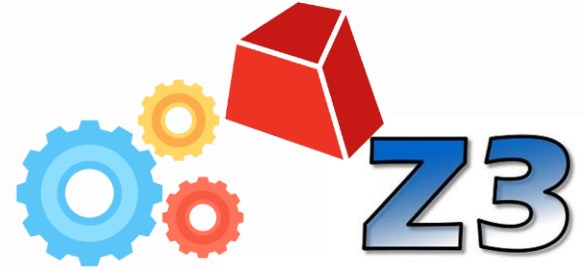
# Decomposing global constraints



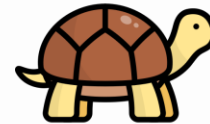
*AllDifferent*( $x, y, z$ )

Modeling system

$x \neq y \wedge x \neq z \wedge y \neq z$



IBM  
CPLEX



Efficient Modeling of  
**Half-reified**  
global constraints



# Reified constraints

“Reification relates the truth value of a constraint to a Boolean variable”

Full-reification:  $b \leftrightarrow \text{AllDifferent}(x, y, z)$

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But... Full reification is hard!

Requires to assert the negation of the constraint if  $b$  is False.

E.g.,  $\neg \text{Circuit}(x, y, z)$

# Half-reified constraints

- Full reification can be "overkill"
  - $b \Leftrightarrow \text{AllDifferent}(x, y, z)$
- Often only need half-reification
  - $b \rightarrow \text{AllDifferent}(x, y, z)$
- Easy to derive half-reified propagators!

## Half Reification and Flattening

Thibaut Feydy<sup>1</sup>, Zoltan Somogyi<sup>1</sup>, and Peter J. Stuckey<sup>1</sup>

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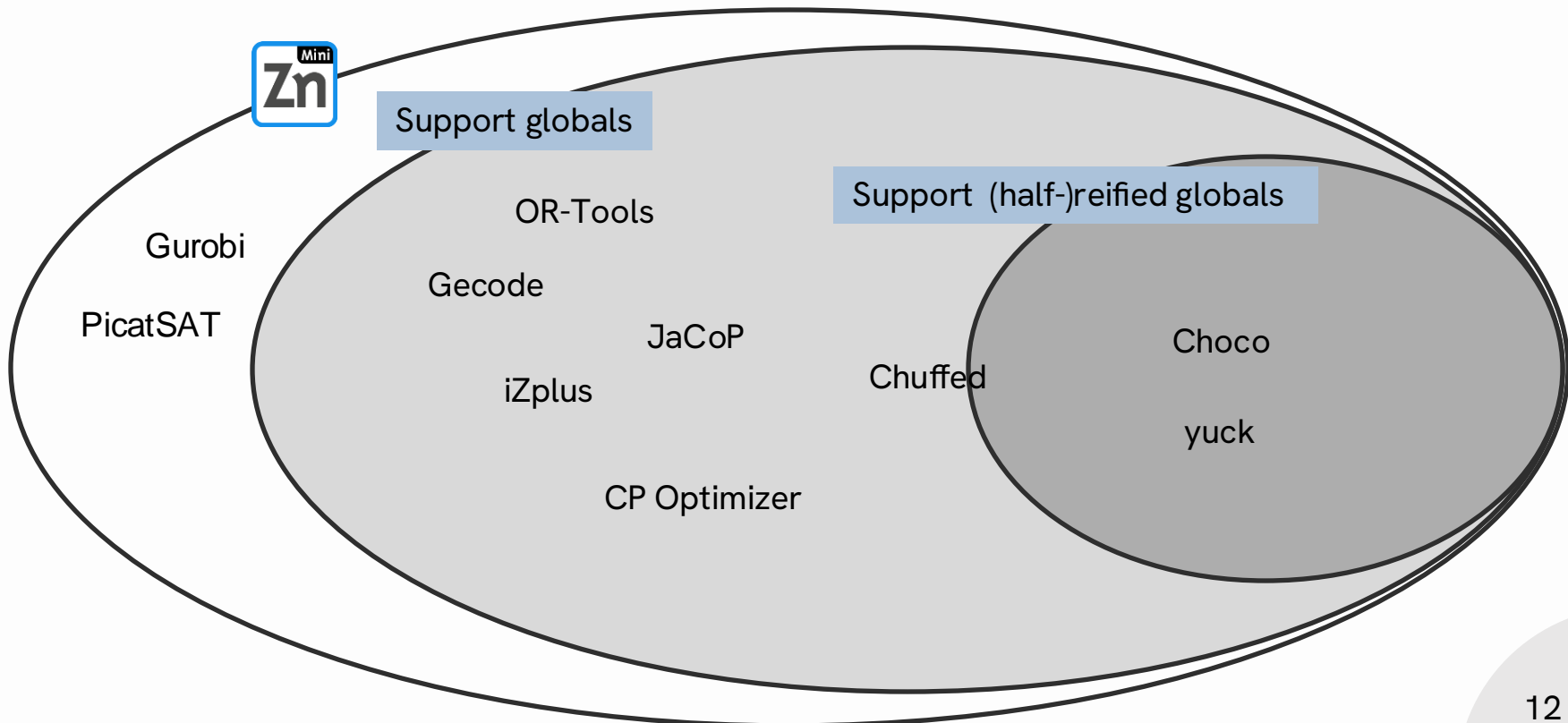
$$\forall v \in \mathcal{V} : f_{b \rightarrow G(\mathcal{V})}(\mathcal{D})[v] = \begin{cases} f_{G(\mathcal{V})}(\mathcal{D})[v] & \text{if } \mathcal{D}[b] = \{true\} \\ \mathcal{D}[v] & \text{otherwise} \end{cases}$$
$$f_{b \rightarrow G(\mathcal{V})}(\mathcal{D})[b] = \begin{cases} \mathcal{D}[b] \setminus \{true\} & \text{if } f_{G(\mathcal{V})}(\mathcal{D}) \text{ is a false domain} \\ \mathcal{D}[b] & \text{otherwise} \end{cases}$$

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# Solver support for (half-)reification



# Decomposing reified global constraints



$b \rightarrow \text{AllDifferent}(x, y, z)$

Modeling system



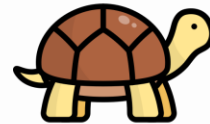
OR-Tools

Gecode JaCoP

iZplus

CP Optimizer

$b \rightarrow (x \neq y \wedge x \neq z \wedge y \neq z)$

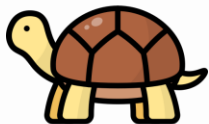


# Why should you care about half-reified global constraints?



# Many applications

- Max-CSP solving
- Incremental solving
- Assumption-based solving
- MUS-computation
- Flattening [by modeling languages]
- ...



$$\begin{array}{l} \text{Maximize } \sum w_c \cdot b_c \\ \text{st. } \quad b_c \rightarrow c \quad \forall c \in C \end{array}$$

**Instead, rewrite with aux vars!**

$b \rightarrow \textit{AllDifferent}(x, y, z)$

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$\text{AllDifferent}(x', y', z') \wedge (b \rightarrow (x = x' \wedge y = y' \wedge z = z'))$



## Instead, rewrite with aux vars!

$$\text{AllDifferent}(x', y', z') \wedge (b \rightarrow (x = x' \wedge y = y' \wedge z = z'))$$

When  $b$  is True:

$$\begin{aligned} &\text{AllDifferent}(x', y', z') \wedge x = x' \wedge y = y' \wedge z = z' \\ &\text{AllDifferent}(x, y, z) \end{aligned}$$

When  $b$  is False:

$$\text{AllDifferent}(x', y', z') \wedge \text{True}$$

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When  $b$  is False:

$$\text{AllDifferent}(x', y', z') \wedge \text{True}$$

+ No decomposition

- Extra variables in model

- Worse propagation compared to native solver support

# Special case: functional constraints

$$b \rightarrow \text{Min}(v, X)$$

$$\text{Min}(v', X) \wedge (b \rightarrow v = v')$$

When  $b$  is True:

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Same propagation strength as native support!

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Same propagation strength as native support!

Also works for full-reification!

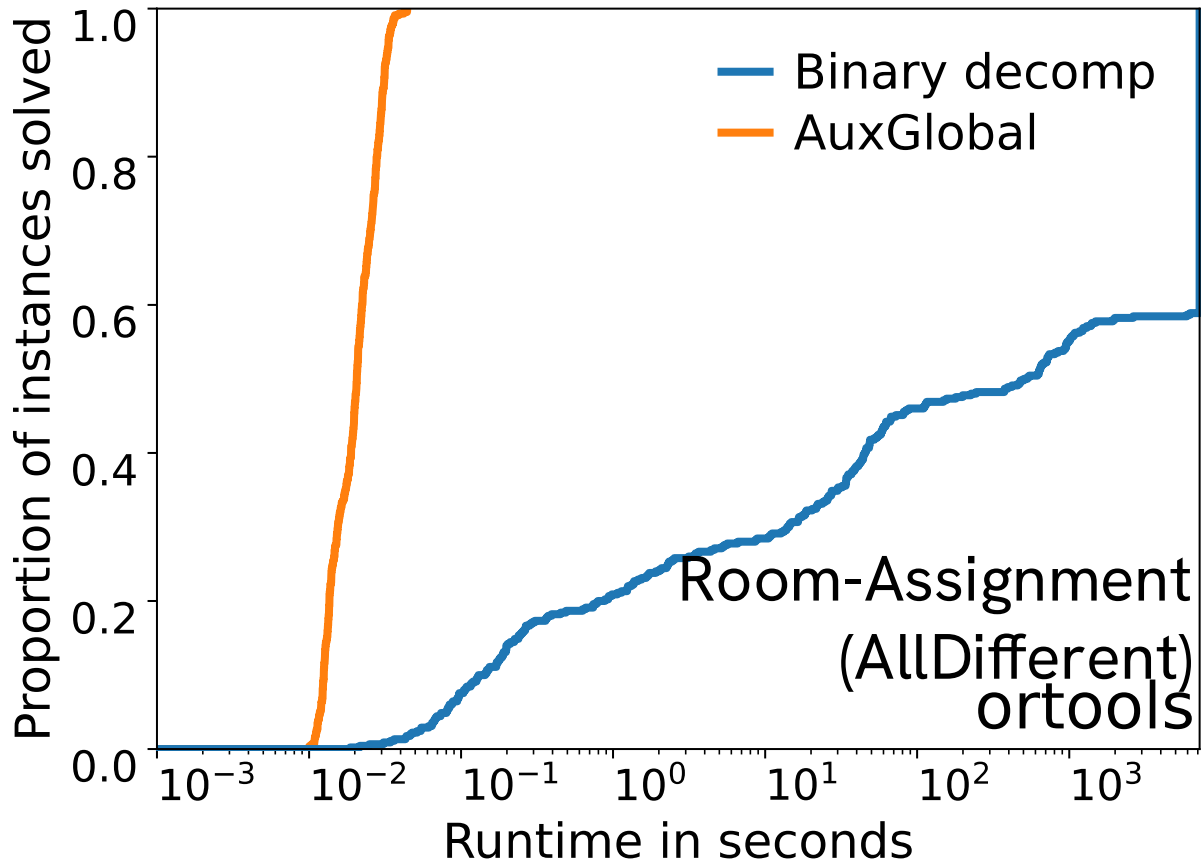
Implemented in most modeling systems!

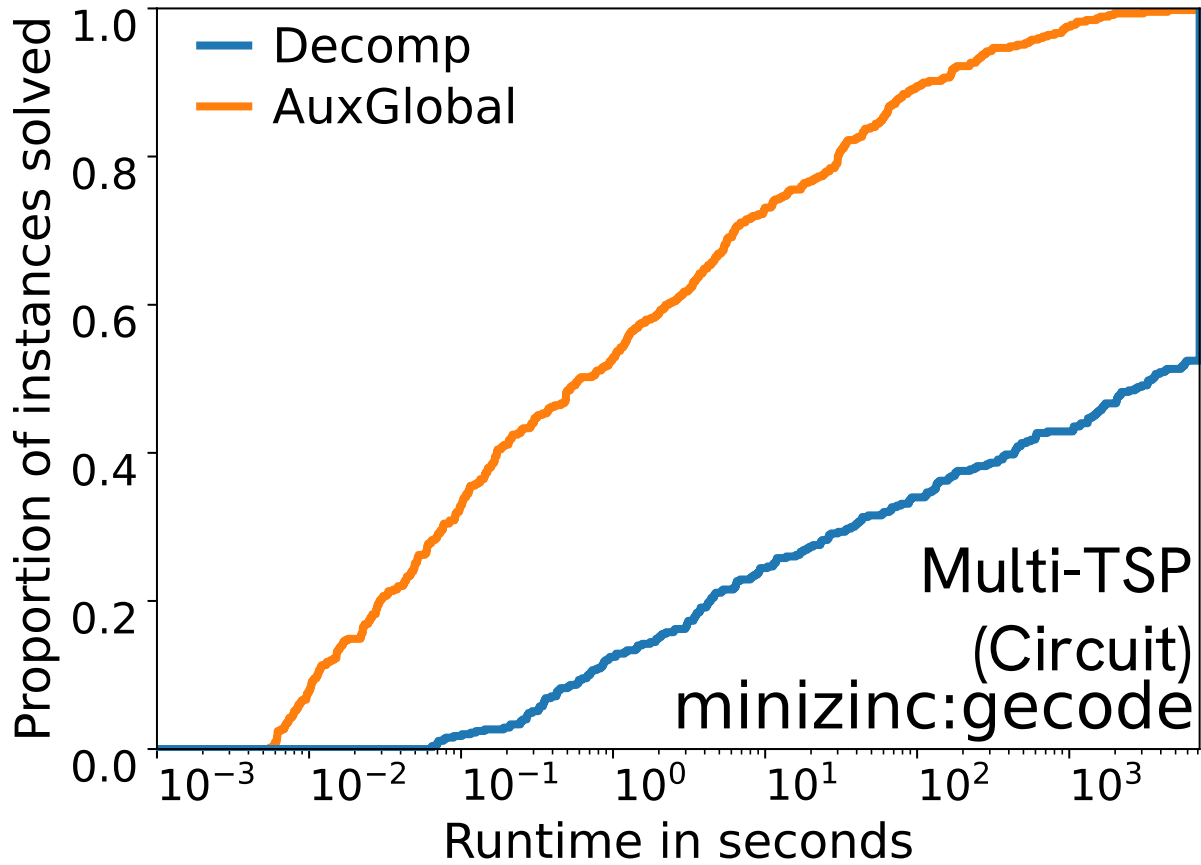
## On the reification of global constraints

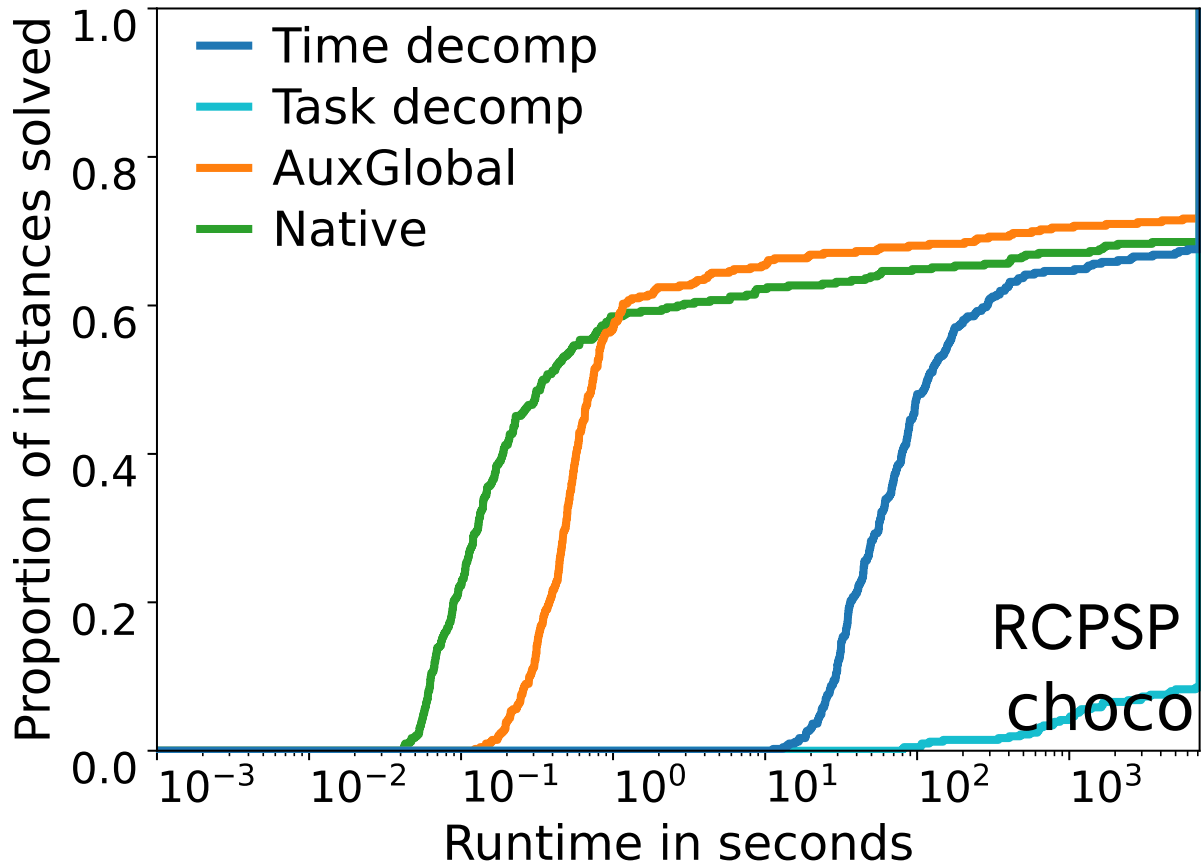
Nicolas Beldiceanu · Mats Carlsson · Pierre Flener ·  
Justin Pearson

# So... Does it work?

- Max-CSP benchmarks with different global constraints:
  - Room-assignment: AllDifferent
  - Multi-TSP: Circuit
  - RCPSP: Cumulative
- Tested on:
  - OR-Tools
  - Gecode
  - Choco
- Through the CPMpy modeling system





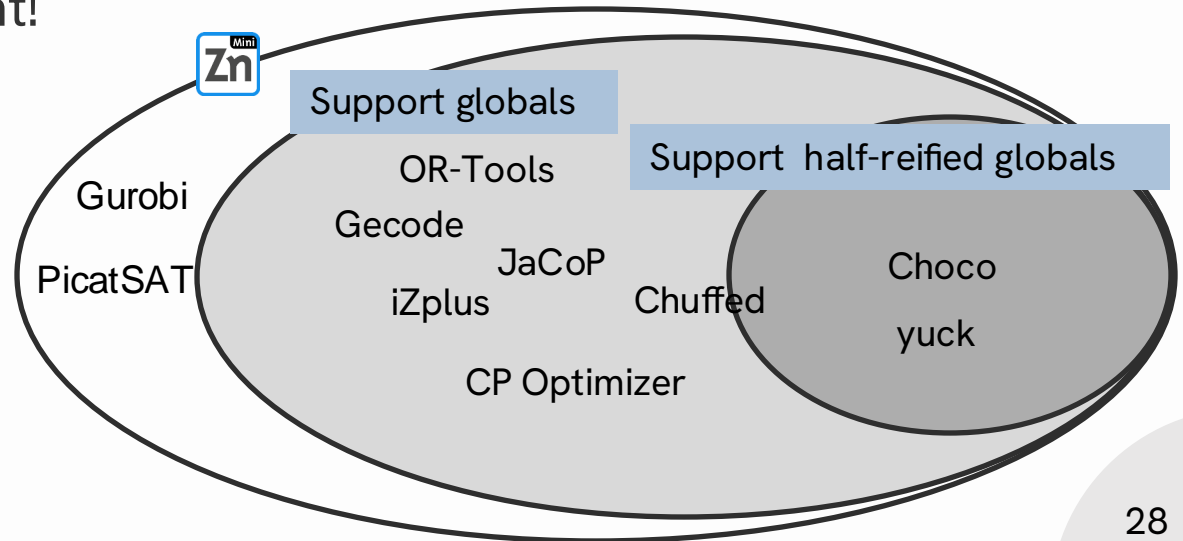


RCPSP  
choco



# Conclusions

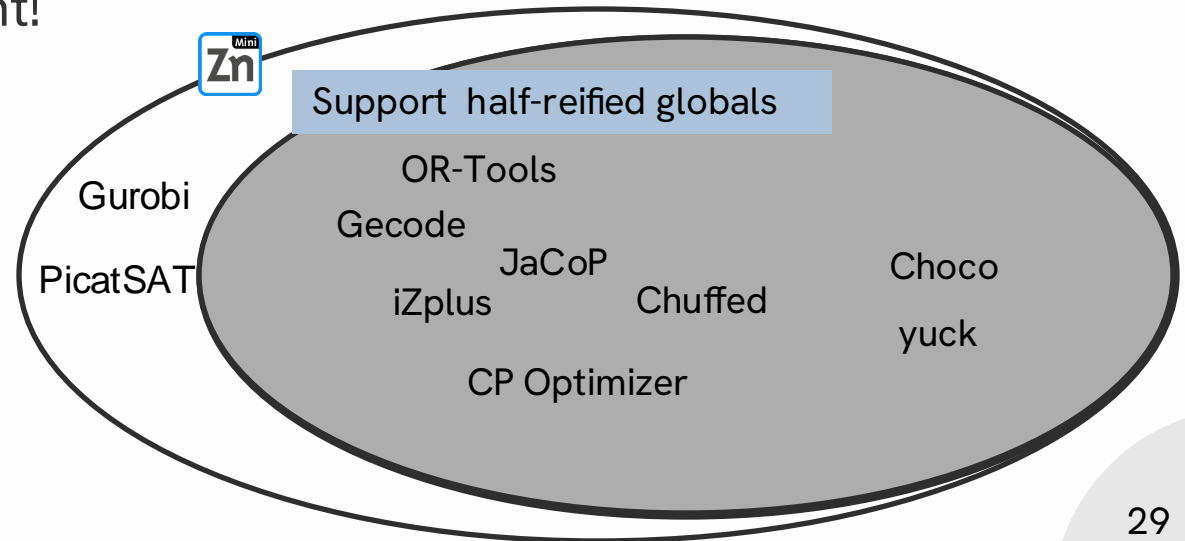
- Reformulation with auxiliary variables is much faster compared to decomposing
- Competitive with solver-native approaches
- Easy to implement!



# Conclusions

- Reformulation with auxiliary variables is much faster compared to decomposing
- Competitive with solver-native approaches
- Easy to implement!

Enable support for half-reified global constraints for ANY CP-solver



# Next steps

- Compare with state-of-the-art Max-CSP solvers
- Evaluate flattening use-cases?
- Minimize auxiliary variables for non-functional constraints

# Questions?